# Large independent sets in triangle-free subcubic graphs: beyond planarity

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Joint work with Jan Goedgebeur and Gwenaël Joret

Université Libre de Bruxelles

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#### G graph on n vertices

Independence number:  $\alpha(G) = \max\{k \mid \exists k \text{ pairwise nonadjacent vertices}\}\$ 

'Subcubic' = maximum degree at most 3

How large is  $\frac{\alpha(G)}{n}$  in subcubic triangle-free graphs G?

## Independence ratio.

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 $\Rightarrow$ 

More feasible goal:

Given a (nice) graph class  $\mathcal{G}$ , find largest  $c \in (0,1]$  such that for all  $G \in \mathcal{G}$ :

$$\frac{\alpha(G)}{n} \geq c.$$

## Planar graphs

Theorem (four-colour theorem, 1976, Appel and Haken)

G planar  $\Rightarrow G$  is 4-colourable.

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*G* planar  $\Rightarrow$  *G* is 4-colourable.

#### Corollary

$$G$$
 planar  $\Rightarrow \alpha(G) \geq \frac{n}{4}$ .

Every known proof uses the four-colour theorem!

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#### Corollary

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What if we also forbid triangles?

# Triangle-free planar

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## Theorem (1985, Jones)

G triangle-free planar  $\Rightarrow \alpha(G) \geq \frac{n+1}{3}$ .

Sharp for infinitely many graphs

Graph class	$\alpha \geq$	Due to
planar	<u>n</u>	4-colour theorem (1976)
planar triangle-free	<u>n+1</u>	Grötszch (1959), Jones (1985)

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Can we do better than  $\frac{1}{3}$ ?

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Can we do better than  $\frac{1}{3}$ ?

Yes, if we assume G is subcubic.

# Subcubic triangle-free

#### Theorem (1979, Statton)

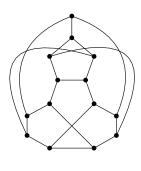
G subcubic triangle-free  $\Rightarrow \alpha(G) \geq \frac{5}{14}n$ .

# Subcubic triangle-free

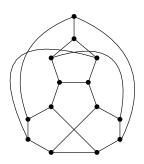
#### Theorem (1979, Statton)

G subcubic triangle-free 
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.

Only two tight examples among connected graphs [Heckman '08].



$$n = 14$$
  $\alpha = 5$ 



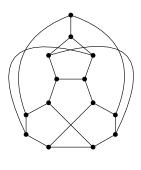
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# Subcubic triangle-free and connected

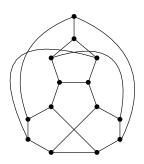
#### Theorem (1995, Fraughnaugh and Locke)

G connected subcubic triangle-free 
$$\Rightarrow \alpha(G) \ge \frac{11}{30}n - \frac{4}{30}$$
.

The same two graphs are the only tight examples:



$$n = 14$$
  $\alpha = 5$ 

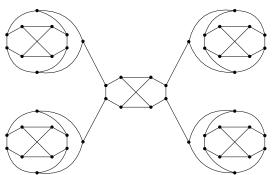


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# Subcubic triangle-free and connected

#### Theorem (1995, Fraughnaugh and Locke)

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Asymptotically tight; infinitely many graphs with  $\alpha = \frac{11}{30}n - \frac{2}{15}$ .

Graph class	$\alpha \geq$	Due to
planar	<u>n</u>	4-colour theorem (1976)
planar triangle-free	<u>n+1</u>	Grötszch (1959), Jones (1985)
subcubic triangle-free	<u>5<i>n</i></u> 14	Statton (1979)
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Can we do better than  $\frac{11}{30}$ ?

Yes, if we additionally assume G is planar.

	1	I .
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planar subcubic triangle-free	3 <i>n</i> 8	Conjectured (1976) by Albertson, Bollóbas and Tucker, proved by Heckman and Thomas (2006)

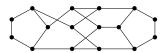
#### Theorem (2006, Heckman and Thomas)

G planar subcubic triangle-free 
$$\Rightarrow \alpha(G) \geq \frac{3}{8}n$$

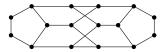
Sharp for infinitely many bad graphs, e.g. :



$$\alpha = 3$$
,  $n = 8$ .







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planar	<u>n</u> 4	4-colour theorem (1976)
planar triangle-free	<u>n+1</u> 3	Grötszch (1959), Jones (1985)
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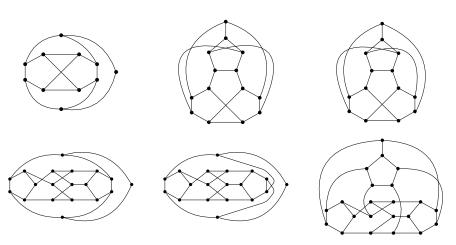
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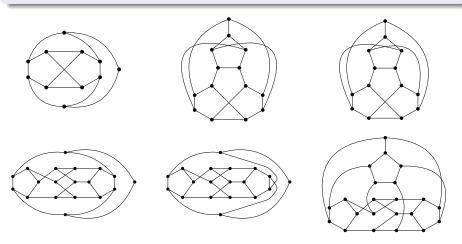
Question: can we relax the planarity condition?

Old conjecture: suffices to forbid six non-planar subgraphs



#### Conjecture (1995, Fraughnaugh and Locke + Bajnok and Brinkmann)

G F-free subcubic triangle-free  $\Rightarrow \alpha(G) \geq \frac{3}{8}n$ .



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Definition: a graph G (on at least three vertices) is called 2-connected if G - v is connected for every vertex v.

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Theorem (2019+, C., Goedgebeur and Joret) Both conjectures are true :-)

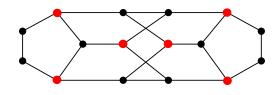
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2-connected subcubic triangle-free	$\frac{3n-2}{8}$	C, Goedgebeur, Joret (2019+)

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Our technical theorem generalizes all of the above results concerning subcubic graphs

Graph class	$\alpha \geq$	Sharp due to
planar	<u>n</u>	$\infty$ many
planar triangle-free	$\frac{n+1}{3}$	$\infty$ many
subcubic triangle-free	<u>5n</u> 14	$\infty$ many (ess. two members of ${\cal F}$ )
connected subcubic triangle-free	$\frac{11n-4}{30}$	two members of ${\cal F}$
planar subcubic triangle-free	3 <i>n</i> 8	$\infty$ many
${\mathcal F}$ -free subcubic triangle-free	3 <u>n</u> 8	∞ many graphs that are 3-connected, girth 5, non-planar and in fact with arbitrarily large clique minors
2-connected subcubic triangle-free	<u>3<i>n</i>−2</u>	three members of ${\mathcal F}$

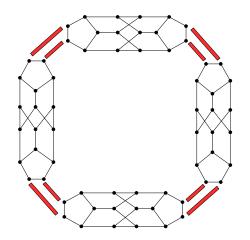
# $\infty$ many sharp examples



$$\alpha = 6 = \frac{3n}{8}$$
.

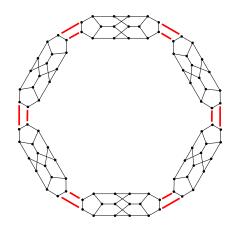
Use as a building block.

# $\infty$ many sharp examples



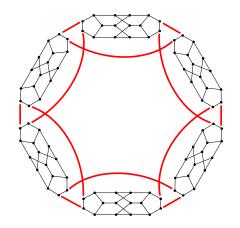
3-connected, non-planar and  $\alpha = \frac{3n}{8}$ .

# $\infty$ many sharp examples



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3-connected, girth 5, non-planar and  $\alpha = \frac{3n}{8}$ .

## Motivation for the technical theorem

## Theorem, C., Goedgebeur, Joret (2019+)

If G is  $\mathcal{F}$ -free subcubic triangle-free, then  $\alpha(G) \geq \frac{3n}{8}$ .

### Enough to show the statement when

- G connected, and
- G critical, meaning  $\alpha(G e) > \alpha(G) \quad \forall e \in E(G)$

A sparsity measure:

$$\mu := \frac{9n_3 + 10n_2 + 11n_1 + 12n_0}{24} = \frac{6n - |E(G)|}{12}$$

where  $n_i :=$  number of vertices of degree i.

Remarks:

$$\mu \geq \frac{3}{8}n$$
 ;

$$\left\lceil \frac{3}{8}n - \frac{1}{12} \right\rceil \ge \frac{3}{8}n$$
 because  $n \in \mathbb{Z}$ .

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Hence, to show  $\alpha \geq \frac{3}{8}n$  it is enough to prove  $\alpha \geq \mu - \frac{1}{12}$ 

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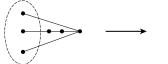
**Problem**: there exist bad graphs with  $\alpha = \mu - \frac{2}{12}$ .

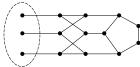
# Definition of bad graphs



is bad and

every 8-augmentation of a bad graph is bad:





**Attempt 2**: Show that  $\alpha \ge \mu - \frac{1}{12}$ , unless *G* is bad.

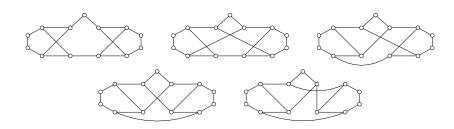
This is true, however . . .

**Problem**: to prove it, we need to consider a slightly stronger statement. This involves dangerous graphs, some of which attain  $\alpha = \mu - \frac{1}{12}$ .

## Dangerous graphs

Technical recursive definition.

(C<sub>5</sub> and the following are the smallest dangerous graphs)



# **Main technical theorem** (CvB-G-J '19+) Suppose

- *G F*-free subcubic and triangle-free
- G connected and critical, and
- G not bad

then 
$$\alpha \geq \mu - \frac{1}{12}$$
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- *G F*-free subcubic and triangle-free
- G connected and critical, and
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then 
$$\alpha \geq \mu - \frac{1}{12}$$
 .

#### If moreover

- G has  $\geq 3$  degree-2 vertices and
- G not dangerous and has no bad subgraph

then 
$$\alpha \geq \mu$$
 .

# Generalization: allowing triangles

Define the refined measure

$$\mu^*(G) := \frac{6n - e - 2t}{12},$$

where e is the number of edges in G and t is the maximum number of vertex-disjoint triangles.

## Theorem (2019+, C, Goedgebeur and Joret)

Let G be a critical connected subcubic graph which is not isomorphic to  $K_4$  or any member of  $\mathcal{F}$ . Then

- $\alpha(G) = \mu^*(G) \frac{2}{12}$  if G is bad or almost bad
- $\alpha(G) \ge \mu^*(G) \frac{1}{12}$  otherwise.

# Plan of the proof

#### G minimum counter-example

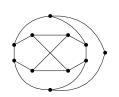
- G almost 3-connected: If X is a 2-cutset then G-X has exactly two components, with one isomorphic to  $K_1$  or  $K_2$
- G has no bad subgraph
- Deal with degree-2 vertices:
  - case where neighbors have both degree 2
  - case where neighbors have both degree 3
  - case where neighbors have degree 2 and 3
- $\rightarrow$  G is cubic and 3-connected
  - G has no 4-cycle
  - G has no 6-cycle
  - G has no dangerous subgraph (in particular, no 5-cycle)

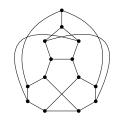
Final argument: Local structure around a shortest even cycle

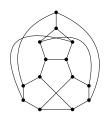
## Open problems

### G $\mathcal{F}$ -free subcubic triangle-free

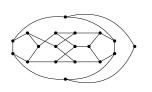
- $\alpha(G) \ge \frac{3n}{8}$  is still best possible if we also forbid four-cycles. What happens for even larger girth?
- Fractional chromatic number  $\chi_f(G)$  at most  $\frac{8}{3}$ ? (cp. conjecture Heckman and Thomas '08)

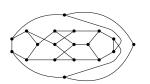


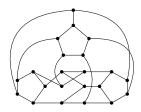




Thank you for your attention!







# An empty slide

# Open problems

**Staton** '79 If *G* subcubic and triangle-free then  $\frac{n}{\alpha} \leq \frac{14}{5}$ 

Recall: 
$$\frac{n}{\alpha} \le \chi_f$$

**Dvořák, Sereni, Volec** '14 (conjectured by Heckman & Thomas '01) If G subcubic and triangle-free then  $\chi_f \leq \frac{14}{5}$ 

Could the upper bound on  $\chi_f$  be improved if we further assume

- G connected, or
- G 2-connected, or
- G planar, or
- G has none of the 6 exceptional graphs as subgraph?

# Plan of the proof; how the graphs of $\mathcal{F}$ emerge.

#### G minimum counter-example

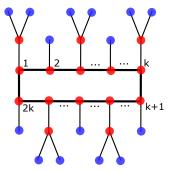
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- G has no bad subgraph.  $\Rightarrow F_{11}, F_{19}^{(1)}, F_{19}^{(2)}$ .
- Deal with degree-2 vertices:
  - case where neighbors have both degree 2.
  - case where neighbors have both degree  $3 \Rightarrow \text{red}$ . to dangerous graphs.
  - case where neighbors have degree 2 and 3  $\Rightarrow$  red. to 8-augmentation.
- $\rightarrow$  G is cubic and 3-connected  $\Rightarrow$  henceforward only need:  $\alpha \ge \mu \frac{1}{12}$ .
  - G has no 4-cycle
  - G has no 6-cycle  $\Rightarrow F_{14}^{(1)}, F_{14}^{(2)}$ .
  - G has no dangerous subgraph (in particular, no 5-cycle)  $\Rightarrow F_{22}$ .

Final argument: Local structure around a shortest even cycle

# Final argument (simplified):

Shortest even cycle  $1, 2, \ldots, 2k$  in G.

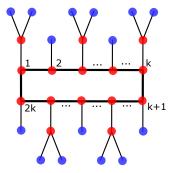
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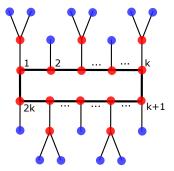


$$\alpha(G) - \alpha(G') \ge k$$
  
=  $3k \cdot \frac{9}{24} - k \cdot \frac{1}{24} = \mu(G) - \mu(G')$ .

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$$\alpha(G) - \alpha(G') \geq k$$

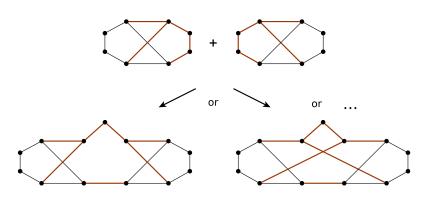
$$= 3k \cdot \frac{9}{24} - k \cdot \frac{1}{24} = \mu(G) - \mu(G').$$

Apply induction to  $\alpha(G') - \mu(G')$  for desired bound on  $\alpha(G) - \mu(G)$ .

# Definition dangerous graphs

### C<sub>5</sub> is dangerous

Join of two bad graphs is dangerous:



$$\alpha = \mu - \frac{1}{12}$$
 if G dangerous.